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A CATEGORICAL GLIMPSE AT THE RECONSTRUCTION OF GEOMETRIES

Dedicated to Karl H. Hofmann, on the occasion of his 60th birthday

ABSTRACT. The author's reconstruction method ['Reconstruction of incidence geometries from groups of automorphisms', Arch. Math. 58 (1992) 621–624] is put in a categorical setting, and generalized to geometries with an arbitrary number of 'types'. The results amount to saying that the reconstruction process involves a pair of adjoint functors, and that the class of those geometries that are images under reconstruction forms a reflective subcategory.

The present paper deals with the problem whether a given geometry (i.e. roughly spoken, an incidence structure together with a group of automorphisms) is determined by the action of the group. After results on flagtransitive groups on incidence structures with two types of objects (see [7] for further historical notes), a first attempt by the author led to results that were quite satisfactory for the case of partial planes (see Proposition (2.7) below), and in particular for stable planes (see [7, 4]). In a discussion about this previous result, K. H. Hofmann urged a categorical point of view. This categorical treatment led to new insights: some of the previous restrictions on the action of the group may be dropped, and the reconstruction process applies to geometries with an arbitrary number of types. Moreover, the categorical point of view helps the groups to their full rights: in fact, this paper presents an equivalence (in a strictly categorical sense) of 'sufficiently homogeneous' geometries and systems of subgroups.

1. A CATEGORY OF INCIDENCE STRUCTURES

(1.1) NOTATION. By Set and Gp we denote the familiar categories of sets and mappings, and of groups and group homomorphisms, respectively. The fact that a morphism (in any category) is monic (epic) shall be stressed by the notation \rightarrow (\rightarrow). An isomorphism shall be denoted by \Rightarrow , which is not to be confused with \Rightarrow (logical implication). A natural number is defined as the set of its predecessors: $n := \{m | m < n\}$. Thus 0 is the empty set, and $1 = \{0\}$. For each mapping $f: X \rightarrow Y$ and each subset $U \subseteq Y$ we write $f^{-}(U) := \{x | f(x) \in U\}$.

We shall consider categories that are not balanced, i.e. a morphism that is

both monic and epic need not be an isomorphism. Hence the following classes of morphisms are of interest:

(1.2) DEFINITION. Consider any category.

(a) A monomorphism m is called *extreme*, if for each factorization $m = h \circ g$ the fact that g is epic implies that g is an isomorphism:

(b) Dually, an epimorphism e is called *extreme*, if for each factorization $e = h \circ g$ the fact that h is monic implies that h is an isomorphism:

It is easy to see from the definition that an extreme monomorphism is an isomorphism if, and only if, it is an epimorphism; and dually. Note that **Set** and **Gp** are balanced categories, where, in particular, each monic and each epic is extreme. We shall denote extreme monics and epics by the symbols c_{p} and $\neg b_{p}$, respectively.

(1.3) DEFINITION. An incidence structure (I, A) consists of a family $A: T \to \mathbf{Set}: t \mapsto A_t$ and a subset $I \subseteq \prod_{t \in T} A_t$. For incidence structures (I, A) and (J, B), a morphism $f = (F, f_T): (I, A) \to (J, B)$ consists of a mapping $f_T: T \to U$ (where U is the domain of B), and a family $F: t \mapsto F_t$ with domain T such that $F_t: A_t \to B_{f_T(t)}$ is a mapping (for each $t \in T$). Finally we require that there exists some mapping $F_I: I \to J$ such that

(A)
$$I \xrightarrow{i} \prod_{t \in T} A_t \xrightarrow{\operatorname{pr}_t} A_t$$
$$\downarrow^{F_t} \qquad \qquad \downarrow^{F_t}$$
$$J \xrightarrow{j} \prod_{u \in U} B_u \xrightarrow{\operatorname{pr}_{f_T(t)}} B_{f_T(t)}$$

is commutative for each $t \in T$, where *i* and *j* denote the inclusion maps, and pr, the projection. Every mapping F_I that meets these requirements shall be called *suited*. These data define a category, which we shall denote by **Inc**. The

following diagram makes clear how the composition $g \circ f$ of morphisms has to be defined (and that this definition works well):

(1.4) INTERPRETATION. The domain T of A may be considered as set of types of geometric objects, while I is an incidence relation in the spirit of G. Pickert's definition* [6: 1.1, p. 2]. In this interpretation, a morphism is a mapping that preserves incidence and equality of type. Note that the 'type mapping' f_T need not be monic nor epic; and that, in general, the mapping F_I is not determined uniquely.

(1.5) LEMMA. Assume that $f = (F, f_T): (I, A) \rightarrow (J, B)$ is a morphism in Inc such that f_T is bijective. Then the following holds:

- (a) The suited mapping F_I is determined uniquely by the family F.
- (b) If F_t is monic for each $t \in T$ then F_t is monic.

Proof. If f_T is a bijection then there exists the product mapping

$$\varphi := \prod_{t \in T} (F_t \circ \mathrm{pr}_t): \prod_{t \in T} A_t \to \prod_{t \in T} B_{f_T(t)}$$

and the following diagram commutes for each suited F_I :

Since *j* is monic, this implies that the suited mapping F_I is unique. This proves assertion (a). If each of the F_t is monic, then φ is monic. Hence F_I is monic, and assertion (b) is established.

^{*}In contrast with a 'symmetric' notion of incidence, cf. [3].

Note that a mapping with the properties of the product mapping (in particular, uniqueness!) need not exist if f_T is not epic* or not monic.

(1.6) **PROPOSITION**. Let $f = (F, f_T): (I, A) \rightarrow (J, B)$ be a morphism in Inc.

- (a) f is epic if, and only if, f_T and F_t for each $t \in T$ (but not necessarily F_I) are epic.
- (b) f is monic if, and only if, f_T and F_t for each $t \in T$ are monic.
- (c) An epimorphism f is extreme if, and only if, every suited F_1 is epic.
- (d) A monomorphism f is extreme if, and only if **,

$$\left(\prod_{t\in T} (F_t \circ \operatorname{pr}_t)\right)^{\leftarrow} \left(\left(\prod_{t\in T} \operatorname{pr}_{f_T(t)}\right)(J)\right) = I.$$

Proof. (i) Assume that f_T is epic, and that F_t is epic for each $t \in T$. We consider morphisms $g = (G, g_U): (J, B) \to (K, C)$ and $h = (H, h_U): (J, B) \to (K, C)$ such that $g \circ f = h \circ f$. This means that $g_U \circ f_T = h_U \circ f_T$ and that $G_{f_T(t)} \circ F_t = H_{f_T(t)} \circ F_t$ for each $t \in T$. The fact that f_T and the F_t are epimorphisms implies that $g_U = h_U$ and that $G_u = H_u$ for each $u \in U = f_T(T)$. Hence f is epic.

(ii) Conversely, assume that either f_T or one of the F_t is not epic. We define a ('2-gon') incidence structure (K, C) by V := 2, $C_0 := 2 =: C_1$, and $K := 2 \times 2$. Obviously, the mappings $g_U \equiv 0$, $G_u \equiv 0$ define a morphism $g = (G, g_U): (J, B) \rightarrow (K, C)$: e.g. $G_J \equiv (0, 0)$ is suited. On the other hand, there exists a morphism $h = (H, h_U)$ such that

$$h_U(u) = \begin{cases} 0 & \text{if } u \in f_T(T) \\ 1 & \text{otherwise} \end{cases}$$
$$H_u(x) = \begin{cases} 0 & \text{if } u = f_T(t) \text{ and } x \in F_t(A_t) \\ 1 & \text{otherwise.} \end{cases}$$

Now $g \circ f = h \circ f$, but g = h only if f_T and each of the F_t is epic. Thus assertion (a) is established.

(iii) Dualizing the arguments of (i), one obtains that f is monic if f_T and each of the F_t are monic.

(iv) Assume that there are two types s, $t \in T$ such that $f_T(s) = f_T(t)$. We consider the incidence structure (K, C) that is given by the settings V := 1,

^{*}For example, consider the embedding of a line into 3-dimensional affine space: one is free to choose any plane that is incident with the line.

^{**}If $f_T = \mathrm{id}_T$ this condition simplifies to the condition that $(\prod_{i \in T} F_i)^- (J) = I$. Note that the inclusion ' \supseteq ' is already a consequence of the fact that F_I is suited.

 $C_0 := 0$, and K := 0. There are morphisms $g, h: (K, C) \to (I, A)$ such that $g_K(0) = s$ and $h_K(0) = t$. Now $f \circ g = f \circ h$, and we infer that f is not monic. If, on the other hand, there is some $t \in T$ such that there are two elements $x, y \in A_t$ with $F_t(x) = F_t(y)$, then V' := 1, $C'_0 := 1$, and K' := 0 define an incidence structure (K', C') that admits morphisms to (I, A) that show that f is not monic. Thus assertion (b) is established.

(v) In order to prove assertion (c), we fix a suited mapping F_I and consider the incidence structure (J', B) that is given by $J' := F_I(I)$. Obviously, the morphism f may be regarded as a morphism from (I, A) to (J', B). Now the diagram

$$\begin{array}{cccc} (I,A) & \stackrel{f}{\longrightarrow} & (J,B) \\ f & & \parallel \\ (J',B) & \stackrel{h}{\longrightarrow} & (J,B) \end{array}$$

commutes, where $h_U := id_U$, and $H_u := id_{B_u}$ for each $u \in U$ (the inclusion $H_{J'}: J' \to J$ is suited). If f is an extreme epimorphism, we conclude that h is an isomorphism, hence J' = J.

(vi) Assume that f is an epimorphism and that every suited F_I is epic. If $f = h \circ g$ for morphisms $g:(I, A) \to (K, C)$ and $h:(K, C) \to (J, B)$ and h is monic, then h is epic (since f is), and we infer that h_K and each of the H_v are bijections. Since $F_I := H_K \circ G_J$ is suited, it is epic, and this holds for H_K , too. Thus H_K is a bijection by Lemma (1.5(b)), and h is an isomorphism. This proves assertion (c).

(vii) For the proof of assertion (d) we introduce the ad hoc notation

$$\begin{aligned} \pi &:= \prod_{t \in T} \operatorname{pr}_{F_{T}(t)} : \prod_{u \in U} B_{u} \to \prod_{t \in T} B_{f_{T}(t)} \\ \varphi &:= \prod_{t \in T} (F_{t} \circ \operatorname{pr}_{t}) : \prod_{t \in T} A_{t} \to \prod_{t \in T} B_{f_{T}(t)} \\ I' &:= \varphi^{\leftarrow}(\pi(J)) = \left(\prod_{t \in T} (F_{t} \circ \operatorname{pr}_{t})\right)^{\leftarrow} \left(\left(\prod_{t \in T} \operatorname{pr}_{f_{T}(t)}\right) (J)\right) \end{aligned}$$

We obtain an incidence structure (I', A). From the commutativity of (A) we infer that $I \subseteq I'$. Now for each $x \in I'$, there exists some $y_x \in J$ such that $\varphi(x) = \pi(y_x)$. Any choice of such a y_x for each $x \in I'$ defines a suited mapping $F_{I'}: x \mapsto y_x$. Hence we have the following commuting diagram:



Here for each $t \in T$, the identity id_{A_t} is denoted by E_t , and $E_I: I \subseteq I'$ is the inclusion.

(viii) If f is an extreme monic, we conclude that $(E, id_T): (I, A) \to (I', A)$ is an isomorphism. This implies that E_I is a bijection (cf. (1.5)), hence I' = I.

(ix) Finally, assume that I = I', and that f is monic. Let $g = (G, g_T): (I, A) \rightarrow (K, C)$ and $h = (H, h_V): (K, C) \rightarrow (J, B)$ be morphisms such that g is epic and $f = h \circ g$. Then g is monic (since f is), hence g_T and each of the G_t are bijections. According to (1.5) there is a unique suited mapping G_I , and G_I is monic. We have the following diagram:

Here $\gamma := \prod_{t \in T} G_t \circ pr_t$. From the relation

$$\varphi^{\leftarrow}(\pi(H_K(K))) \subseteq \varphi^{\leftarrow}(\pi(J)) = I'$$

we infer that

$$K \subseteq \gamma(\varphi^{\leftarrow}(\pi(H_K(K)))) \subseteq \gamma(I').$$

Since $\gamma(I') = \gamma(I) = G_I(I)$ and k is monic, we conclude that G_I is epic. Hence g is an isomorphism, and assertion (d) is proved.

Later on, we shall be interested mainly in morphisms (F, f_T) that preserve types (in the sense that $f_T = id_T$, or at least is a bijection). In this case, extreme epimorphisms are easier to recognize:

(1.7) LEMMA. Assume that $f = (F, f_T): (I, A) \rightarrow (J, B)$ is an epimorphism in Inc, and that f_T is a bijection. Then f is extreme if, and only if, there exists a suited F_I that is epic.

Proof. This follows immediately from (1.5(a)) and (1.6(c)).

(1.8) REMARK. Similar arguments may be used to show that extreme monics (epics) in Inc are embeddings (identifications)*.

(1.9) EXAMPLES. (a) In descriptive geometry, one is particularly interested in extreme monomorphisms: e.g. a drawing of a polyhedron may be interpreted as a morphism from a (finite) incidence structure to the real affine plane. Of course one tries to avoid that the drawing 'shows' incidences that do not occur in the original object.

(b) Closure theorems give examples of situations where one is interested in monomorphisms that are not extreme. For example, the validity of Desargues' theorem may be phrased in the following way: Let (I, A) be the incidence structure that is obtained by 'drawing two perspective triangles, the three points where corresponding sides meet, and a line that contains exactly two of these three points'. Then Desargues' theorem is valid in (J, B) if there is no monomorphism from (I, A) to (J, B) that is extreme.

(c) Examples of morphisms that change types are given by dualities. Further examples are the embeddings of rank 2 residues in diagram geometries, cf. [3].

2. GEOMETRIES, SKETCHES AND PICTURES

We turn to geometries in a sense inspired by F. Klein:

(2.1) DEFINITION. Assume that (I, A) is an incidence structure, and let Γ be a group. A geometry (γ, I, A) is given by a family γ (with domain T) of group actions $\gamma_t: \Gamma \times A_t \to A_t$ such that there exists an action $\gamma_I: \Gamma \times I \to I$ making the following diagram commutative:

id- x pr

(B

Note that γ_I is uniquely determined, cf. (1.5).

A morphism $(\varphi, f) = (\varphi, F, f_T): (\gamma, I, A) \rightarrow (\delta, J, B)$ consists of a morphism $\varphi: \Gamma \to \Delta$ (in **Gp**) and a morphism $f = (F, f_T): (I, A) \to (J, B)$ such that commutes for some suited F_I . These data define a category, which we shall denote by **Geo**. There are the forgetful functors

 $\mathbf{U}_{\mathbf{Inc}}$: **Geo** \rightarrow **Inc**: $(\gamma, I, A) \mapsto (I, A)$ and $\mathbf{U}_{\mathbf{Gp}}$: **Geo** \rightarrow **Gp**: $(\gamma, I, A) \mapsto \Gamma$.

*Here we use the terms 'embedding' and 'identification' in the sense of 'Einbettung' and 'Identifizierung', as defined in [2, 4.8.1, 4.10].



(2.2) **PROPOSITION.** There are initial and terminal objects (but no null objects^{*}) in the categories **Inc** and **Geo**:

- (i) There exists exactly one object of Inc with no types at all; in this case, the family A (and hence also I) is empty. The corresponding incidence structure (0,0) is initial in Inc. Fixing the (cardinality of the) set T of types, one obtains a subcategory Inc_T of Inc. In Inc_T, the initial object is given by (0, (0)_{teT}), where (0)_{teT}: T → Set: t → 0.
- (ii) The (unique) action of the initial 1 ∈ Gp on the empty set gives rise to the initials ℑ and ℑ_T in Geo and Geo_T, respectively, where U_{Inc}(ℑ) = (0,0) and U_{Inc}(ℑ_T) = (0, (0)_{t∈T}).
- (t) The setting $T = A_0 = I = 1$ defines a terminal object $(1, (1)) \in \text{Inc}$, while I = 1 and $A \equiv 1$ gives the terminal object $(1^T, (1)_{t \in T}) \in \text{Inc}_T$.
- (tt) The (unique) action of 1 on 1 gives terminals \mathfrak{T} and \mathfrak{T}_T in Geo and Geo_T, respectively, where $\mathbf{U}_{Inc}(\mathfrak{T}) = (1, (1))$ and $\mathbf{U}_{Inc}(\mathfrak{T}_T) = (1^T, (1)_{t \in T})$.

The proof of these statements is straightforward and left to the reader.

The initial $1 \in \mathbf{Gp}$ acts on each incidence structure (K, C) in a unique way; we shall denote this action by (v, K, C). In particular, we obtain that the functor $\mathbf{U_{Inc}}$ is full. (This is in contrast to familiar forgetful functors, e.g. $\mathbf{Gp} \rightarrow \mathbf{Set.}$) Note that for each group Ψ there are unique actions $\Psi \times 0 \rightarrow 0$ and $\Psi \times 1 \rightarrow 1$. This implies that there are families ι and τ of actions such that $(\iota, 0, 0)$ and $(\tau, 1, (1))$ are geometries. We shall use these objects in the proof of the following statement.

(2.3) LEMMA. The functors U_{Inc} and U_{Gp} preserve monics and epics.

Proof. (i) Assume that f is not epic. Then there are morphisms $g, h: (J, B) \rightarrow (K, C)$ such that $g \circ f = h \circ f$ but $g \neq h$. Now the unique morphism $\omega: (\delta, J, B) \rightarrow \mathfrak{T}$ factors through morphisms (α, g) and (β, h) from (δ, J, B) to

^{*}If one requires the sets I, T, and A, to be non-empty, then (1, (1)) and $(1^T, (1)_{teT})$ are null objects in Inc and Inc_T, respectively; while \mathfrak{T} and \mathfrak{T}_T are null objects in Geo and Geo_T, respectively.

 (ν, K, C) . Now $\alpha \neq \beta$ implies that $(\alpha, g) \neq (\beta, h)$, but $(\alpha, g) \circ (\varphi, f) = (\beta, h) \circ (\varphi, f)$, hence (φ, f) is not epic. This shows that U_{Inc} preserves epics. Dual arguments (with \Im instead of \Im) show that U_{Inc} preserves monics.

(ii) Assume that φ is not epic. Then there are morphisms $\alpha, \beta: \Delta \to \Psi$ such that $\alpha \circ \varphi = \beta \circ \varphi$ but $\alpha \neq \beta$. Now the unique morphism $\omega: (\delta, J, B) \to \mathfrak{T}$ factors through morphisms (α, g) and (β, h) from (δ, J, B) to $(\tau, 1, (1))$. Now $\alpha \neq \beta$ implies that $(\alpha, g) \neq (\beta, h)$, but $(\alpha, g) \circ (\varphi, f) = (\beta, h) \circ (\varphi, f)$, hence (φ, f) is not epic. This shows that U_{Gp} preserves epics. Dual arguments (with $(\iota, 0, 0)$ instead of $(\tau, 1, (1))$) show that U_{Gp} preserves monics.

Note, however, that neither U_{Inc} nor U_{Gp} reflects monics or epics.

(2.4) **PROPOSITION**. Let $(\varphi, f): (\gamma, I, A) \rightarrow (\delta, J, B)$ be a morphism in **Geo**.

- (a) (φ, f) is epic if, and only if, φ and f are epic.
- (b) (φ, f) is monic if, and only if, φ and f are monic.
- (c) An epimorphism (φ, f) is extreme if, and only if, f is extreme.
- (d) A monomorphism (φ, f) is extreme if, and only if, f is extreme.

Proof. According to (2.3), the fact that (φ, f) is epic (monic) implies that both φ and f are epic (monic). Since the composition of morphisms in **Geo** is defined by $(\beta, h) \circ (\alpha, g) := (\beta \circ \alpha, h \circ g)$, the product functor $U_{Gp} \times U_{Inc}$ reflects monics and epics. Thus assertions (a) and (b) are proved, and (c), (d) follow.

Given a geometry $\mathfrak{G} = (\gamma, I, A)$, one may try to recover information about \mathfrak{G} from the family γ of actions. Of course, this attempt will yield interesting results only if this family is 'sufficiently strong'. Our next definition is a precision of 'strength':

(2.5) DEFINITIONS. (a) Let $\mathfrak{G} = (\gamma, I, A)$ be a geometry. A family $L: T \to \mathbf{Set}$ is said to *sketch* \mathfrak{G} , if the following conditions hold:

- For each $t \in T$, the set L_t is a set of representatives for the orbit decomposition of A_t under the action γ_t .
- The product $L_I := \prod_{t \in T} L_t$ is contained in *I*, and forms a set of representatives for the orbit decomposition of *I* under the action γ_I .

If L sketches \mathfrak{G} , then (L, γ, I, A) is called a sketched geometry. A morphism

$$(\varphi, f): (L, \gamma, I, A) \to (M, \delta, J, B)$$

of sketched geometries is a morphism $(\varphi, f): (\gamma, I, A) \to (\delta, J, B)$ in **Geo** such that $F_t(L_t) \subseteq M_t$ for each $t \in T$ (then $F_I(L_I) \subseteq M_J$ for each suited F_I). Thus we obtain a category **SGeo**.

(b) Let Γ be a group, and denote the set of all subgroups of Γ by $\sigma(\Gamma)$. A

sketch (Γ, \mathscr{L}) is given by a family $\mathscr{L}: T \to 2^{\sigma(\Gamma)}: t \mapsto \mathscr{L}_t$ (i.e. \mathscr{L}_t is a set of subgroups of Γ). A morphism $(\varphi, \Phi, f_T): (\Gamma, \mathscr{L}) \to (\Delta, \mathscr{M})$ consists of a morphism $\varphi: \Gamma \to \Delta$ (in **Gp**), a mapping $f_T: T \to U$ (where U is the domain of \mathscr{M}), and a family $\Phi: t \mapsto \Phi_t$ with domain T such that

 $-\Phi_t: \mathscr{L}_t \to \mathscr{M}_{f_T(t)}$ is a mapping for each $t \in T$.

- For each $\Lambda \in \mathscr{L}_t$, the inclusion $\varphi(\Lambda) \leq \Phi_t(\Lambda)$ holds.

These data define a category Sk of sketches.

(2.6) **PROPOSITION.** The family Z, where $Z_t = \{1\}$ for each t, sketches the terminal \mathfrak{T} resp. \mathfrak{T}_T . Thus one obtains null objects in **SGeo**.

Again, the (straightforward) proof is left to the reader.

The existence of a family L that sketches a given geometry (γ, I, A) is rather a strong condition:

(2.7) **PROPOSITION.** Let (I, A) be a partial plane* with point space A_p and line space A_1 , and assume that there are families γ and L such that $(L, \gamma, I, A) \in \mathbf{SGeo}$. Then either L_p or L_l has exactly one element.

Proof. Assume that there are points x, $y \in L_p$ and lines v, $w \in L_l$. Then

 $\{(x, v), (x, w), (y, v), (y, w)\} \subseteq I,$

hence either x = y or v = w by the definition of a partial plane.

In particular, this means that each sketched finite linear space is point homogeneous, cf. [1, Th. 2.1]. Outside the realm of partial planes, one can easily cook up examples with non-trivial sets L_t :

(2.8) EXAMPLE. Assume that the points p_n $(1 \le n \le 4)$ form the vertices of a tetrahedron, and denote the sides of this tetrahedron by s_n in such a way that p_n lies on s_m if $n \ne m$. The transpositions (12) and (34) generate a subgroup Γ of Sym₄ that acts in the obvious way on the vertices and on the sides of the tetrahedron. Now the sets $L_p := \{p_1, p_3\}$ and $L_s := \{s_2, s_4\}$ sketch the given geometry.

There is a rather obvious way to associate a sketch with each sketched geometry $\mathfrak{G} = (L, \gamma, I, A)$, namely to define $\mathscr{L}_t := \{\Gamma_r^t | r \in L_t\}$ where Γ_r^t denotes the stabilizer of r with respect to the action γ_t . Then $\mathbf{S}(\mathfrak{G}) := (\Gamma, \mathscr{L})$ is a sketch. For the reversed direction, from **Sk** to **SGeo**, we have to work a little harder:

^{*}In the sense of M. Hall [4]: i.e. there are two types ('points' and 'lines'): $T = \{p, l\}$, and an incidence relation $I \subseteq A_p \times A_l$ such that for any two points $x, y \in A_p$ there is at most one line $w \in A_l$ with $\{(x, w), (y, w)\} \subseteq I$, and dually.

(2.9) DEFINITION. Let $\mathfrak{S} = (\Gamma, \mathscr{L})$ be a sketch. For each $t \in T$, write

$$P_t := P_t(\mathfrak{S}) := \bigcup_{\Lambda \in \mathscr{L}_t} \frac{\Gamma}{\Lambda},$$

and let λ_t be the action of Γ on P_t via left multiplication. With

$$P_I := P_I(\mathfrak{S}) := \{ (\alpha \Lambda_t)_{t \in T} | \alpha \in \Gamma, \Lambda_t \in \mathscr{L}_t \},\$$

and the action λ_I of Γ on P_I via (simultaneous) left multiplication, we obtain a geometry (λ, P_I, P) :

Obviously, the family \mathscr{L} sketches (λ, P_I, P) . (Note that P_I is just the union of the orbits of the elements of $\prod_{t \in T} \mathscr{L}_t$ under the action λ_I .) We shall call $(\mathscr{L}, \lambda, P_I, P)$ the *picture* of (Γ, \mathscr{L}) and denote it by $\mathbf{P}(\Gamma, \mathscr{L})$.

3. **RECONSTRUCTION**

The categorical view of the reconstruction of sketched geometries amounts to the proof that the reconstruction process involves a pair of adjoint functors (with object mappings S and P, respectively), and that the class of those geometries that are images under reconstruction forms a reflective subcategory.

Straightforward verification shows:

(3.1) LEMMA. For each morphism (φ, F, f_T) : $(L, \gamma, I, A) \rightarrow (M, \delta, J, B)$ define

$$\mathbf{S}(\varphi, F, f_T) := (\varphi, \Phi, f_T),$$

where Φ_t is given by $\Phi_t(\Gamma_r^t) := \Delta_{F_t(r)}^{f_T(t)}$. Then S: SGeo \rightarrow Sk is a functor.

A first step towards the understanding of the connection between S and P is the following trivial observation:

(3.2) LEMMA. $\mathbf{S} \circ \mathbf{P} = \mathrm{id}_{\mathbf{Sk}}$.

The study of $\mathbf{P} \circ \mathbf{S}$ requires some more effort.

(3.3) THEOREM. For each sketched geometry $\mathfrak{G} = (L, \gamma, I, A)$, there is an extreme epimorphism $\eta_{\mathfrak{G}} = (\mathrm{id}_{\Gamma}, \mathrm{H}, \mathrm{id}_{T}): \mathfrak{G} \to \mathbf{P}(\mathbf{S}(\mathfrak{G}))$. The epimorphism $\eta_{\mathfrak{G}}$ is

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an isomorphism if, and only if, the following holds: If $r, s \in L_t$, the equality $\Gamma_r^t = \Gamma_s^t$ implies that r = s.

Proof. For each $x \in A_t$, there is a (unique) representative $r \in L_t$ and a (usually not unique) element $\alpha \in \Gamma$ such that $x = \gamma_t(\alpha, r)$. If $\gamma_t(\alpha, r) = \gamma_t(\beta, s)$ for some $(\beta, s) \in \Gamma \times L_t$, then $\alpha^{-1}\beta \in \Gamma_t^r$ and r = s. This implies that the mapping

$$H_t: A_t \to P_t: \gamma_t(\alpha, r) \mapsto \alpha \Gamma_r^t$$

is well defined, and that it is injective if, and only if, for any two representatives $r, s \in L_t$ the stabilizers Γ_r^t and Γ_s^t are different. (Note, however, that for different types $t, u \in T$, there may hold an equation $\Gamma_r^t = \Gamma_s^u$ without disturbing the injectivity of any of the H_t , cf. Example (2.8).) Similar arguments show that

$$H_I: I \to P_I: \gamma_I(\alpha, (r_t)_{t \in T}) \mapsto (\alpha \Gamma_{r_t}^t)_{t \in T}$$

is a well-defined mapping. Obviously, the family H defines an extreme epimorphism $\eta_{\mathfrak{G}} = (\mathrm{id}_{\Gamma}, \mathrm{H}, \mathrm{id}_{T})$. The assertion that $\eta_{\mathfrak{G}}$ is an isomorphism is equivalent to the assertion that $\eta_{\mathfrak{G}}$ is monic, hence equivalent to the assertion that for each $t \in T$ the mapping H_t is monic.

(3.4) THEOREM. $\langle \mathbf{S}, \mathbf{P}, \psi \rangle$: SGeo \rightarrow Sk is an adjunction with unit $\eta = (\mathfrak{G} \mapsto \eta_{\mathfrak{G}})$ and counit $\varepsilon = (\mathfrak{S} \mapsto \mathrm{id}_{\mathfrak{S}})$, where

$$(\psi_{\mathfrak{G},\mathfrak{S}})^{-1}$$
: SGeo($\mathfrak{G}, \mathbf{P}\mathfrak{S}$) \rightarrow Sk(S $\mathfrak{G}, \mathfrak{S}$): $(\varphi, F, f_T) \mapsto$ S(φ, F, f_T).

In particular, we have that $\mathbf{P}: \mathbf{Sk} \to \mathbf{SGeo}$ is a functor, and that η is a natural transformation from $\mathrm{id}_{\mathbf{SGeo}}$ to $\mathbf{P} \circ \mathbf{S}$.

Proof. Since $\varepsilon_{(\Gamma,\mathscr{L})} = \mathrm{id}_{(\Gamma,\mathscr{L})}$ is a universal arrow from S to (Γ, L) , this follows from [5, Th. IV.1.2(iv)].

Since each component $\varepsilon_{\mathfrak{S}}$ of the counit ε is an isomorphism, we obtain [5, Th. IV.3.1]:

(3.5) COROLLARY. The functor **P** is full and faithful.

(3.6) THEOREM. The image of Sk under P in SGeo forms a reflective subcategory.

Proof. Using (3.2) we obtain that $\mathbf{P} \circ \mathbf{S} \circ \mathbf{P} = \mathbf{P}$, hence $\mathbf{P} \circ \mathbf{S}|_{\mathbf{P(Sk)}} = id_{\mathbf{P(Sk)}}$. Now [5, Th. IV.1.2(v)] says that $\mathbf{P} \circ \mathbf{S}$ is a left adjoint for the inclusion functor.

Quite often, one wishes to identify a certain object in an incidence structure with the set of objects that are incident with it. Generalizing notions like 'point row' or 'line pencil', we introduce *residues* in an incidence structure (I, A) with type set T by the following procedure: Fix an element r of type u,

consider all objects that are 'in some flag together with r', and introduce an incidence structure with type set $T \setminus \{u\}$ on these objects. To be formal:

(3.7) DEFINITION. Let (I, A) be an incidence structure, and fix some type $u \in T$ and an element $r \in A_u$. Writing $\hat{A}_t := \operatorname{pr}_t(I \cap \operatorname{pr}_u^-(r))$ for each $t \in T \setminus \{u\}$ defines a family \hat{A} of sets with domain $T \setminus \{u\}$. With

$$\widehat{I} := \{ (\mathrm{pr}_t(x))_{t \in T \setminus \{u\}} \mid x \in I \cap \mathrm{pr}_u^{\leftarrow}(r) \} \subseteq \prod_{t \in T \setminus \{u\}} \widehat{A}_t$$

we obtain an incidence structure (\hat{I}, \hat{A}) . To stress the choice of $u \in T$ and $r \in A_u$, we write $\operatorname{Res}_r^u := (\hat{I}, \hat{A})$.

(3.8) **PROPOSITION**. Assume that (L, γ, I, A) is a sketched geometry, and fix $r \in L_u$ for some type $u \in T$. Then $(\hat{L}, \hat{\gamma}, \hat{I}, \hat{A})$ is a sketched geometry, where $(\hat{I}, \hat{A}) = \operatorname{Res}_r^u$ as in (3.7), and

$$\hat{\gamma}_t := \gamma_t|_{\Gamma^u_r \times \hat{A}_t}, \quad \hat{L} := L|_{T \setminus \{u\}}.$$

Proof. By definition of Γ_r^u , we have that $\hat{\gamma}_t(\Gamma_r^u \times \hat{A}_t) \subseteq \hat{A}_t$ for each $t \in T \setminus \{u\}$. This implies that $(\hat{\gamma}, \hat{I}, \hat{A})$ is a geometry. In order to show that \hat{L} sketches $(\hat{\gamma}, \hat{I}, \hat{A})$, we observe that each element $x \in I$ is of the form $x = \gamma_I(\alpha, (r_t)_{t \in T}) = (\gamma_t(\alpha, r_t))_{t \in T}$ for some $\alpha \in \Gamma$ and unique representatives $r_t \in L_t$. Now $\gamma_u(r_u) = r$ implies that $r_u = r$, hence $\alpha \in \Gamma_r^u$ if, and only if, the element x belongs to $I \cap \mathrm{pr}_u^-(r)$. This yields that for each $t \in T \setminus \{u\}$, the set \hat{L}_t is a set of representatives for the action $\hat{\gamma}_t$ of Γ_r^u on \hat{A}_t , and that $\Pi_{t \in T \setminus \{u\}} \hat{L}_t$ forms a set of representatives for the action

$$\hat{\gamma}_{\widehat{I}} = ((\alpha, (x_t)_{t \in T \setminus \{u\}}) \mapsto (\hat{\gamma}_t(\alpha, x_t))_{t \in T \setminus \{u\}}).$$

(3.9) COROLLARY. Assume that $\mathfrak{G} = (L, \gamma, I, A)$ is a sketched geometry. Then for each type $u \in T$, the equality $\Gamma_r^u = \Gamma_t^u$ for $r, t \in L_u$ implies that $\operatorname{Res}_r^u = \operatorname{Res}_t^u$.

Proof. This follows easily from the fact that, for each $t \in T$, there holds the equation

$$\operatorname{pr}_t(I \cap \operatorname{pr}_u^{\leftarrow}(r)) = \gamma_t(\Gamma_r^u \times L_t).$$

In view of (3.9), it seems reasonable to consider the full subcategory **RGeo** of **SGeo**, consisting of those sketched geometries (L, γ, I, A) such that the following condition holds:

(R) For every $t \in T$ and $r, s \in L_t$, the equality $\operatorname{Res}_r^t = \operatorname{Res}_s^t$ implies that r = s.

In fact, it is easy to find examples that show that condition (R) does not imply that for any two objects $x, y \in A_t$ we have the implication $\operatorname{Res}_{x}^{t} = \operatorname{Res}_{y}^{t} \Rightarrow x = y$. Though it seems somewhat arbitrarily, condition (R) singles out an important subcategory:

(3.10) THEOREM. The categories **RGeo** and **Sk** are equivalent (via η).

Proof. From (R) we infer that for $r, s \in L_t$, the equation $\Gamma_r^t = \Gamma_s^t$ implies that r = s. According to (3.3), this means that $\eta_{\mathfrak{G}}$ is an isomorphism for each object \mathfrak{G} of **RGeo**.

(3.11) FINAL REMARK. In our present treatment, the situation that was studied in [7] occurs as a special case, namely: The geometries under consideration have two types ('points' and 'lines'), and are subject to three conditions (R1), (R2), and (R3). Condition (R2) says that one has a sketched geometry, with the additional property (R1) that the actions on the points is transitive (i.e. a property that is quite natural in the realm of linear spaces, cf. (2.7)). Condition (R3) says that stabilizers of different representatives are different. Thus it makes sure that the reconstruction process yields an isomorphism. Our result (3.10) above implies that (R3) in [7] may be replaced by condition (R), without affecting the results of [7]. In particular, the validity of (R3) in the examples [7, 7] holds a fortiori.

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